# **Stochastic Spectral and Conjugate Descent Methods** Dmitry Kovalev<sup>1,2</sup> Eduard Gorbunov<sup>2</sup> Elnur Gasanov<sup>1,2</sup> Peter Richtárik<sup>1,2,3</sup>

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#### 1. Introduction

Consider the optimization problem

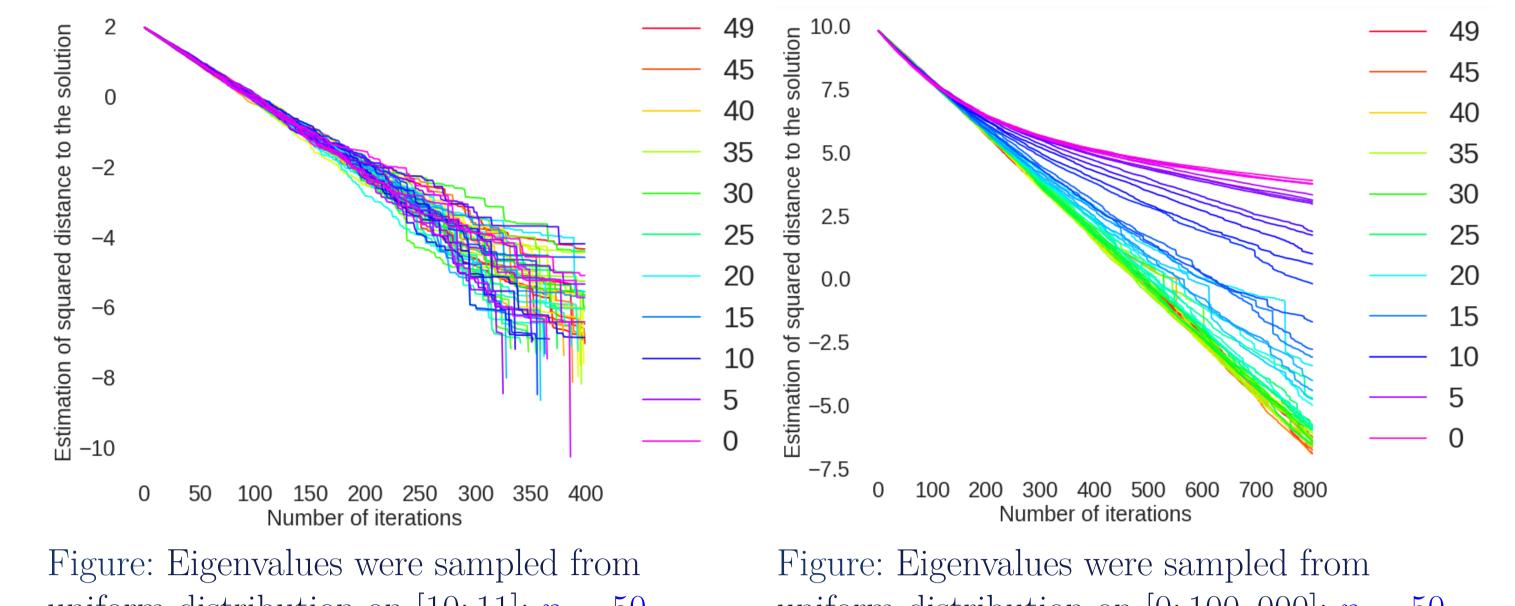
 $\min_{x \in \mathbb{R}^n} f(x) := \frac{1}{2} x^\top \mathbf{A} x - b^\top x,$ 

where **A** is an  $n \times n$  symmetric positive definite matrix. The problem has a unique solution:  $x_* = \mathbf{A}^{-1}b$ . We are interested in the case when n is huge (millions, billions). Note that f is (strongly) convex and quadratic.

## 2. Algorithm: Stochastic Descent

The state-of-the-art methods for convex optimization in huge dimensions are randomized coordinate descent (RCD) methods. We now describe a method

#### 5. Numerical Experiments



which includes RCD as a special case: stochastic descent (SD). SD is a special case of the **sketch-and-project** method developed in [1].

Algorithm 1 [1, 3] (Stochastic Descent).

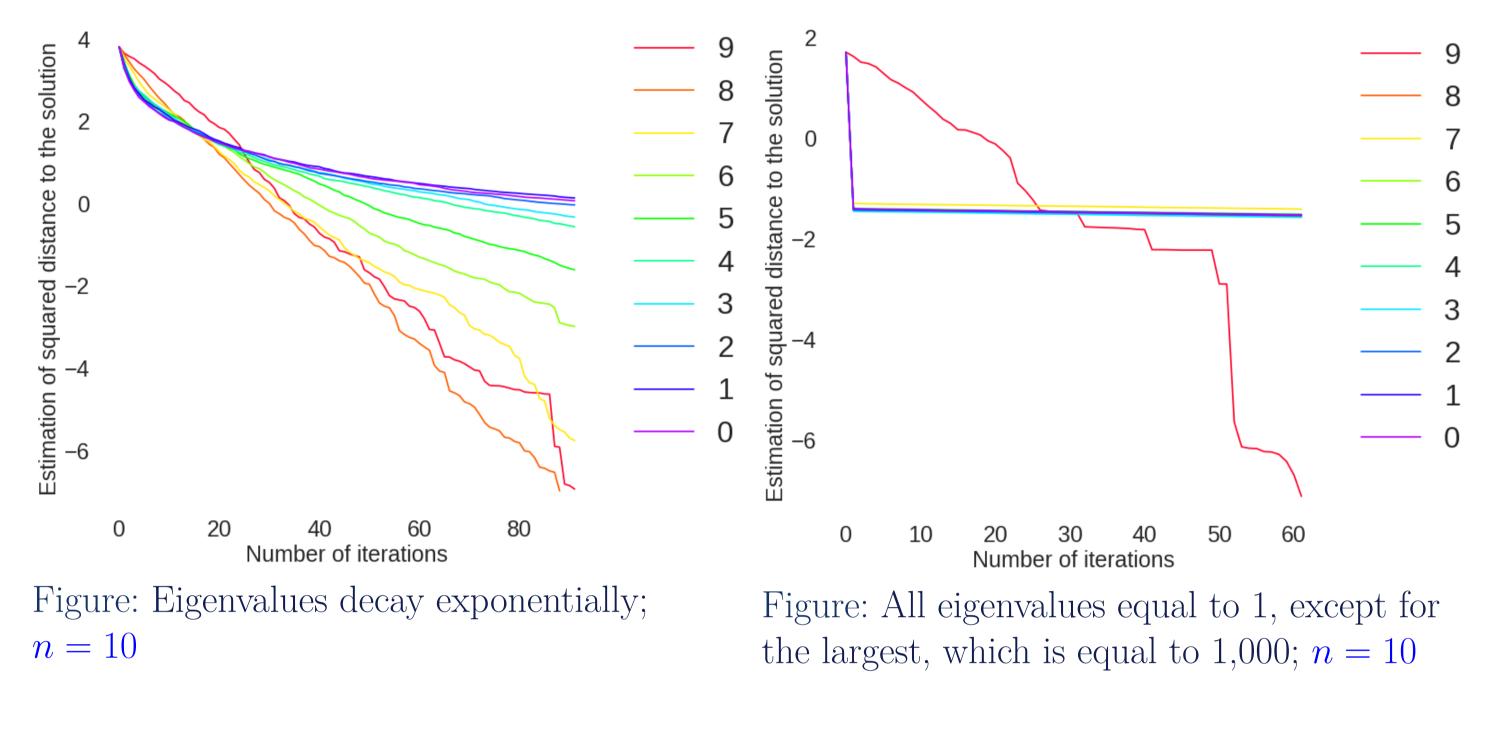
**Parameter:** some distribution  $\mathcal{D}$  over vectors in  $\mathbb{R}^n$ **Initialization:** Choose  $x_0 \in \mathbb{R}^n$ for  $t = 0, 1, 2 \dots do$ Draw a fresh sample  $s_t$  from  $\mathcal{D}$  $x_{t+1} \leftarrow x_t - \frac{s_t^{\top}(\mathbf{A}x_t - b)}{s_t^{\top}\mathbf{A}s_t}s_t$ end for

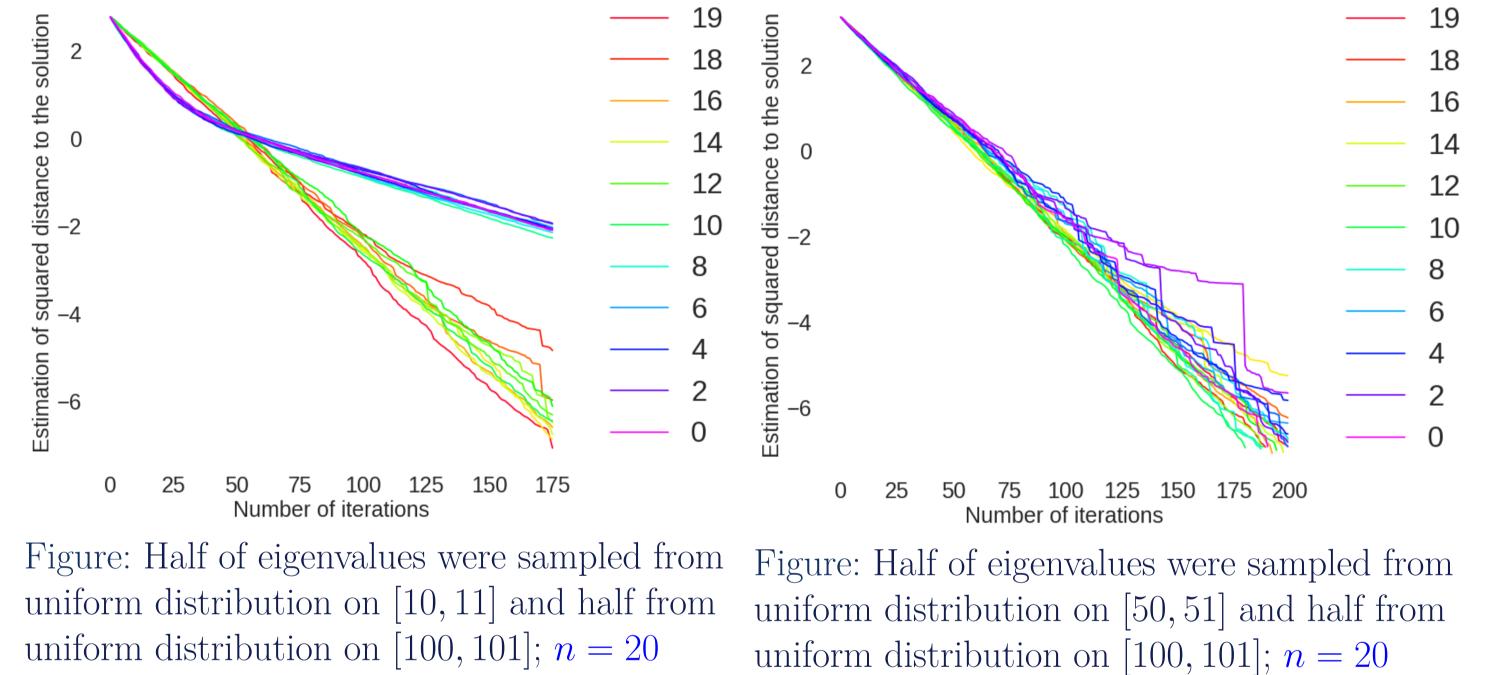
RCD is obtained as a special case by letting  $\mathcal{D}$  be a distribution over unit coordinate (i.e., basis) vectors in  $\mathbb{R}^n$ :  $\{e_1, e_2, \cdots, e_n\}$ :

> $\Leftrightarrow$   $s_t = e_i$  with probability  $p_i > 0$ .  $s_t \sim \mathcal{D}$

**Theorem 1** [1, 3]. Algorithm 1 converges linearly in expectation as  $(1 - \rho_{\max})^t \|x_0 - x_*\|_{\mathbf{A}}^2 \le \mathbb{E}_{s \sim \mathcal{D}}[\|x_t - x_*\|_{\mathbf{A}}^2] \le (1 - \rho_{\min})^t \|x_0 - x_*\|_{\mathbf{A}}^2,$ where  $\|x\|_{\mathbf{A}} = (x^{\top} \mathbf{A} x)^{1/2}$ ,  $\mathbf{W} := \mathbb{E}_{s \sim \mathcal{D}} \left[ \frac{\mathbf{A}^{1/2} s s^{\top} \mathbf{A}^{1/2}}{s^{\top} \mathbf{A} s} \right]$ ,  $\rho_{\max} = \lambda_{\max}(\mathbf{W})$ ,  $\rho_{\min} = \lambda_{\min}(\mathbf{W})$ . Moreover,  $0 < \rho_{\min} \leq 1/n$  and  $\rho_{\max} \leq 1$ .

#### 3. Research Question





RCD with probabilities  $p_i = \mathbf{A}_{ii}/\mathrm{Tr}(\mathbf{A})$  satisfies:  $\rho_{\min} = \lambda_1/\mathrm{Tr}(\mathbf{A})$ , where  $\lambda_1$ is the smallest eigenvalue of A. When  $\rho_{\min}$  is small, RCD is slow. Can we modify RCD by utilizing some spectral information, if known, so that the rate gets improved?

### 4. New Algorithm

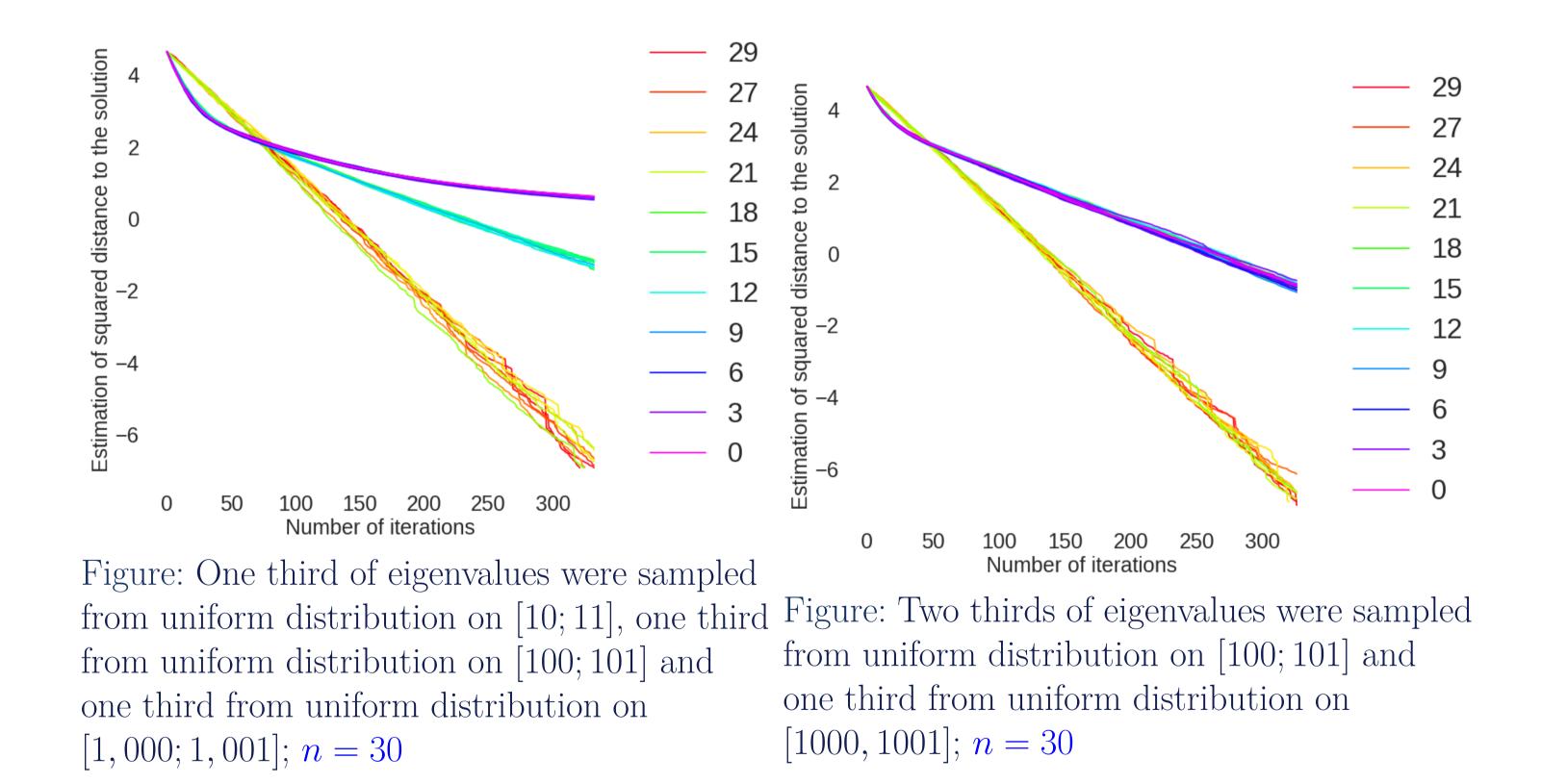
Let  $\mathbf{A} = \sum_{i=1}^{n} \lambda_i u_i u_i^{\mathsf{T}}$  be the eigenvalue decomposition of  $\mathbf{A}$ , with  $0 < \lambda_1 \leq \lambda_2 \leq \lambda_2$  $\cdots \leq \lambda_n$  being the eigenvalues, and  $u_1, \ldots, u_n$  the eigenvectors.

Algorithm 2 [2] (Stochastic Spectral Coordinate Descent). **Parameter:** Choose  $k \in \{0, \ldots, n-1\}$ ; set  $C_k = k\lambda_{k+1} + \sum_{i=k+1}^n \lambda_i$ Run Algorithm 1 with the following distribution  $\mathcal{D}$ :

 $s_t = \begin{cases} e_i & \text{with probability } p_i = \frac{\mathbf{A}_{ii}}{C_k}, \quad i = 1, 2, \dots, n \\ u_i & \text{with probability } p_{n+i} = \frac{\lambda_{k+1} - \lambda_i}{C_k}, \quad i = 1, 2, \dots, k. \end{cases}$ 

Note that for k = 0, Algorithm 2 reduces to RCD.

**Theorem 2.** For every  $n \ge 2$ , Algorithm 2 has the rate



Moreover, the rate improves as k grows, and interpolates between the RCD rate  $\lambda_1/\text{Tr}(\mathbf{A})$  for k = 0, and the optimal rate 1/n for k = n - 1:

$$\frac{\lambda_1}{\operatorname{Tr}(\mathbf{A})} = \frac{\lambda_1}{C_0} \leq \cdots \leq \frac{\lambda_{k+1}}{C_k} \leq \cdots \leq \frac{\lambda_{n-1}}{C_{n-2}} \leq \frac{\lambda_n}{C_{n-1}} = \frac{1}{n}.$$

The total work of Algorithm 2 depends on k:

 $Work(\mathcal{D}) :=$ cost of 1 iteration number of iterations till  $\epsilon$ -solution

k	$P(\mathcal{D})$	$C(\mathcal{D})$	$I(\mathcal{D})$
0	O(n)	O(n)	$\frac{\mathrm{Tr}(\mathbf{A})}{\lambda_1}\ln(1/\epsilon)$
0 < k < n - 1	computation of $\lambda_i$ for $i = 1, 2, \dots, k+1$ computation of $u_i$ for $i = 1, 2, \dots, k$	O(n)	$\frac{C_k}{\lambda_{k+1}}\ln(1/\epsilon)$
n - 1	computation of $\lambda_i$ for $i = 1, 2,, n$ computation of $u_i$ for $i = 1, 2,, n - 1$	O(n)	$n \ln(1/\epsilon)$

#### 6. Bibliography

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