Stochastic Spectral and Conjugate Descent Methods Dmitry Kovalev^{1,2} Eduard Gorbunov² Elnur Gasanov^{1,2} Peter Richtárik^{1,2,3}

¹King Abdullah University of Science and Technology (KAUST), Kingdom of Saudi Arabia

²Moscow Institute of Physics and Technology (MIPT), Russia ³ University of Edinburgh, United Kingdom

1. Introduction

Consider the optimization problem

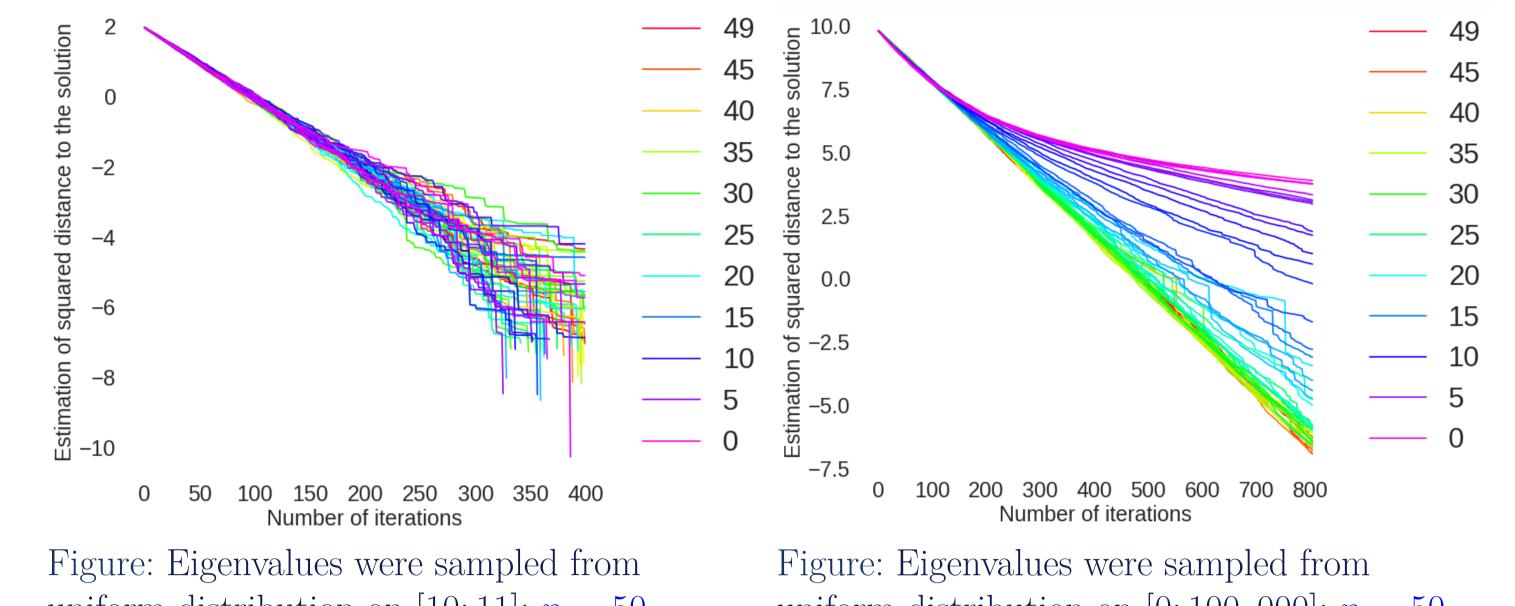
 $\min_{x \in \mathbb{R}^n} f(x) := \frac{1}{2} x^\top \mathbf{A} x - b^\top x,$

where **A** is an $n \times n$ symmetric positive definite matrix. The problem has a unique solution: $x_* = \mathbf{A}^{-1}b$. We are interested in the case when n is huge (millions, billions). Note that f is (strongly) convex and quadratic.

2. Algorithm: Stochastic Descent

The state-of-the-art methods for convex optimization in huge dimensions are randomized coordinate descent (RCD) methods. We now describe a method

5. Numerical Experiments



which includes RCD as a special case: stochastic descent (SD). SD is a special case of the **sketch-and-project** method developed in [1].

Algorithm 1 [1, 3] (Stochastic Descent).

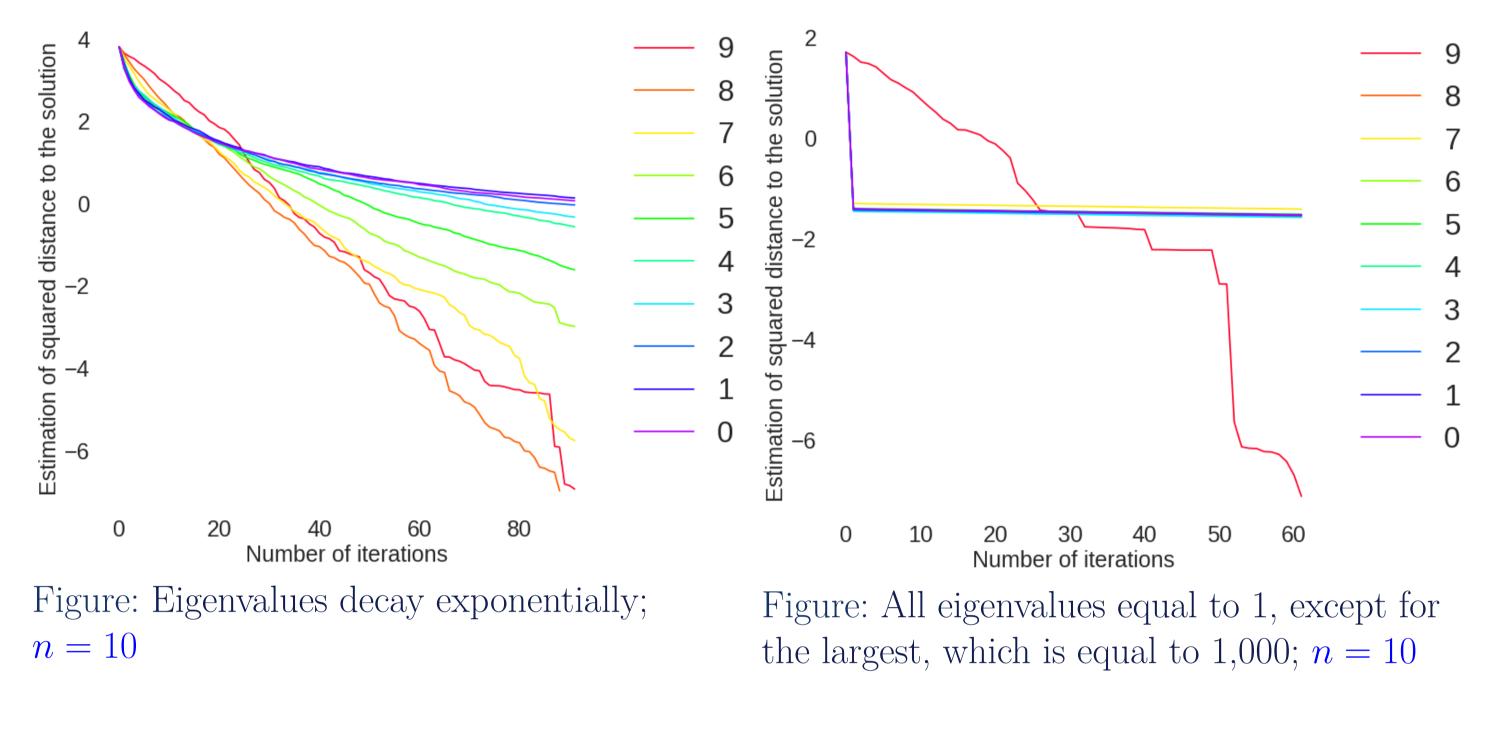
Parameter: some distribution \mathcal{D} over vectors in \mathbb{R}^n **Initialization:** Choose $x_0 \in \mathbb{R}^n$ for $t = 0, 1, 2 \dots do$ Draw a fresh sample s_t from \mathcal{D} $x_{t+1} \leftarrow x_t - \frac{s_t^{\top}(\mathbf{A}x_t - b)}{s_t^{\top}\mathbf{A}s_t}s_t$ end for

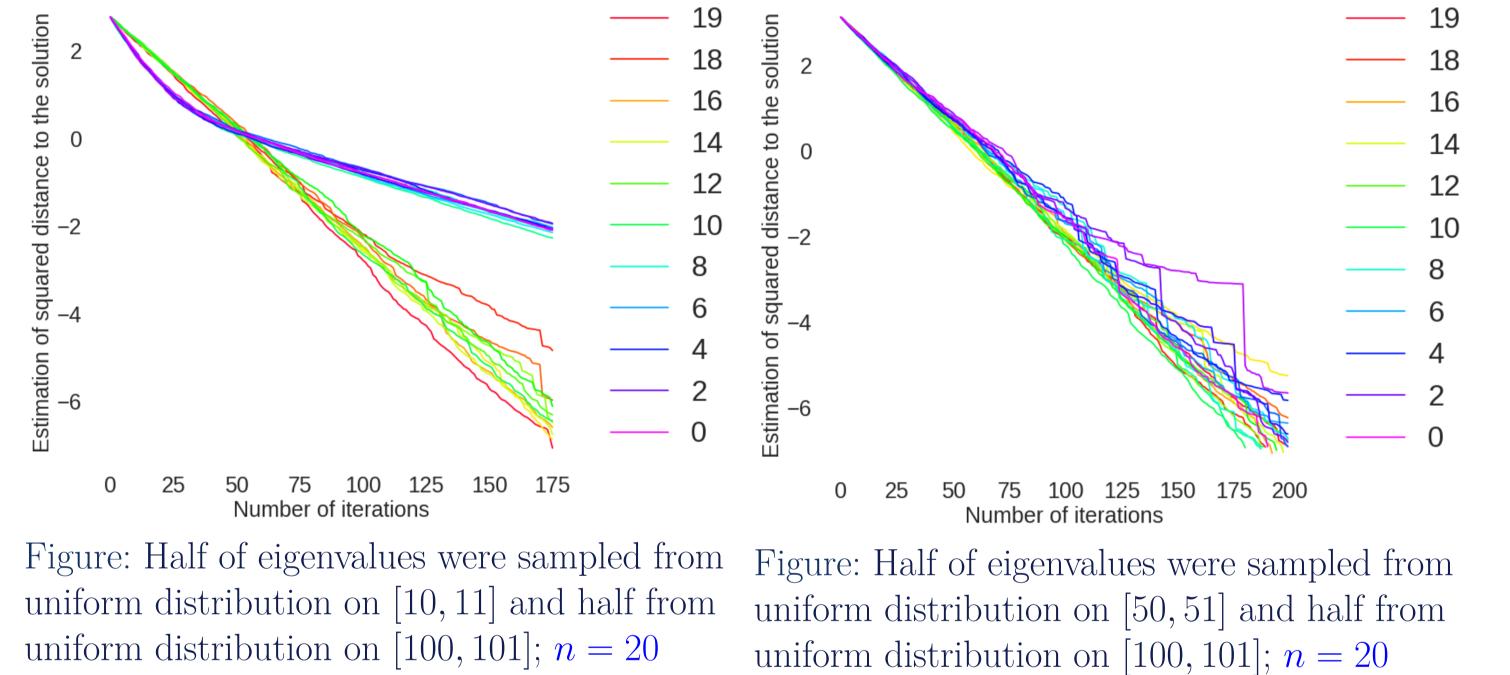
RCD is obtained as a special case by letting \mathcal{D} be a distribution over unit coordinate (i.e., basis) vectors in \mathbb{R}^n : $\{e_1, e_2, \cdots, e_n\}$:

> \Leftrightarrow $s_t = e_i$ with probability $p_i > 0$. $s_t \sim \mathcal{D}$

Theorem 1 [1, 3]. Algorithm 1 converges linearly in expectation as $(1 - \rho_{\max})^t \|x_0 - x_*\|_{\mathbf{A}}^2 \le \mathbb{E}_{s \sim \mathcal{D}}[\|x_t - x_*\|_{\mathbf{A}}^2] \le (1 - \rho_{\min})^t \|x_0 - x_*\|_{\mathbf{A}}^2,$ where $\|x\|_{\mathbf{A}} = (x^{\top} \mathbf{A} x)^{1/2}$, $\mathbf{W} := \mathbb{E}_{s \sim \mathcal{D}} \left[\frac{\mathbf{A}^{1/2} s s^{\top} \mathbf{A}^{1/2}}{s^{\top} \mathbf{A} s} \right]$, $\rho_{\max} = \lambda_{\max}(\mathbf{W})$, $\rho_{\min} = \lambda_{\min}(\mathbf{W})$. Moreover, $0 < \rho_{\min} \leq 1/n$ and $\rho_{\max} \leq 1$.

3. Research Question





RCD with probabilities $p_i = \mathbf{A}_{ii}/\mathrm{Tr}(\mathbf{A})$ satisfies: $\rho_{\min} = \lambda_1/\mathrm{Tr}(\mathbf{A})$, where λ_1 is the smallest eigenvalue of A. When ρ_{\min} is small, RCD is slow. Can we modify RCD by utilizing some spectral information, if known, so that the rate gets improved?

4. New Algorithm

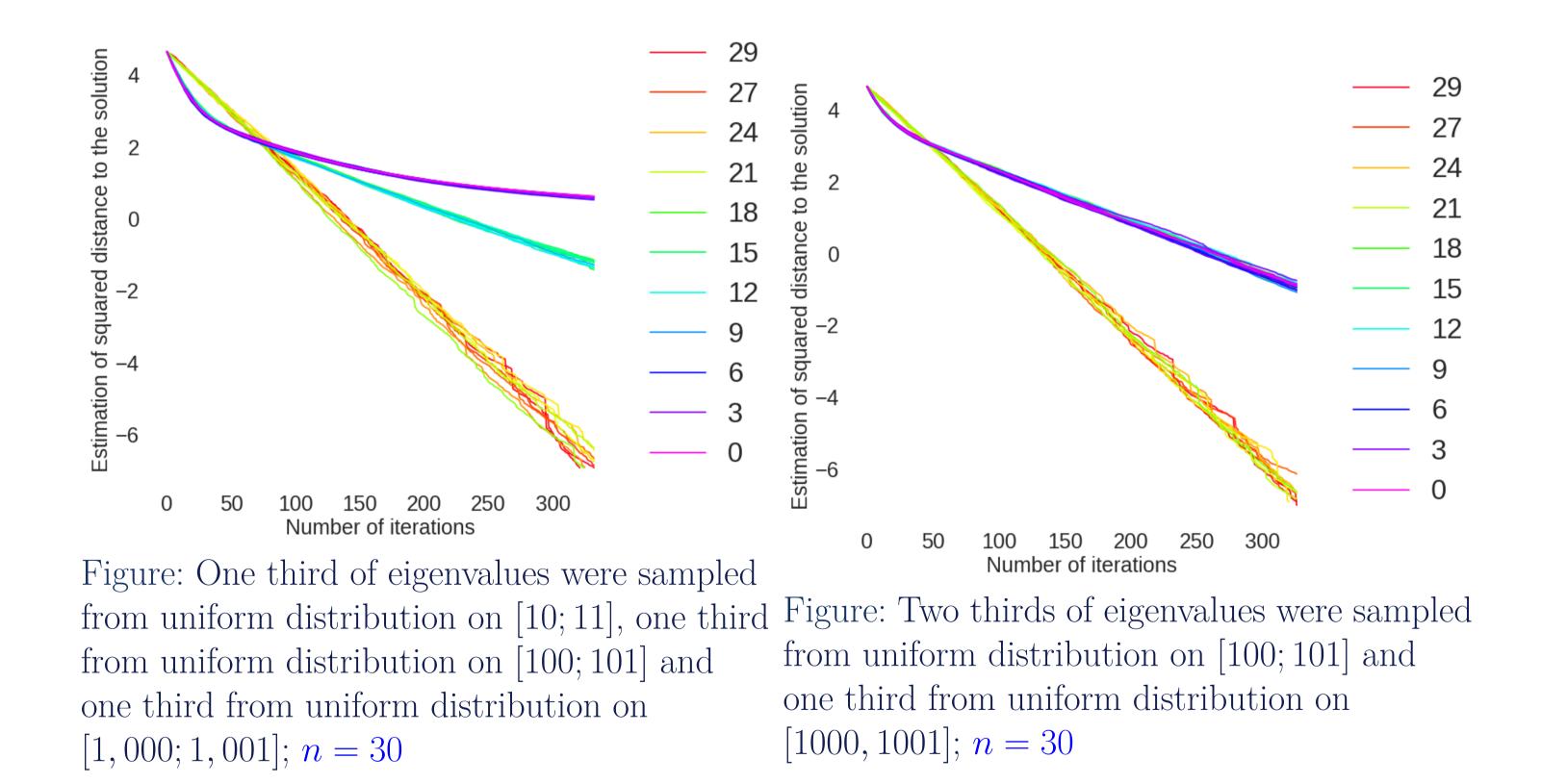
Let $\mathbf{A} = \sum_{i=1}^{n} \lambda_i u_i u_i^{\mathsf{T}}$ be the eigenvalue decomposition of \mathbf{A} , with $0 < \lambda_1 \leq \lambda_2 \leq \lambda_2$ $\cdots \leq \lambda_n$ being the eigenvalues, and u_1, \ldots, u_n the eigenvectors.

Algorithm 2 [2] (Stochastic Spectral Coordinate Descent). **Parameter:** Choose $k \in \{0, \ldots, n-1\}$; set $C_k = k\lambda_{k+1} + \sum_{i=k+1}^n \lambda_i$ Run Algorithm 1 with the following distribution \mathcal{D} :

 $s_t = \begin{cases} e_i & \text{with probability } p_i = \frac{\mathbf{A}_{ii}}{C_k}, \quad i = 1, 2, \dots, n \\ u_i & \text{with probability } p_{n+i} = \frac{\lambda_{k+1} - \lambda_i}{C_k}, \quad i = 1, 2, \dots, k. \end{cases}$

Note that for k = 0, Algorithm 2 reduces to RCD.

Theorem 2. For every $n \ge 2$, Algorithm 2 has the rate



Moreover, the rate improves as k grows, and interpolates between the RCD rate $\lambda_1/\text{Tr}(\mathbf{A})$ for k = 0, and the optimal rate 1/n for k = n - 1:

$$\frac{\lambda_1}{\operatorname{Tr}(\mathbf{A})} = \frac{\lambda_1}{C_0} \leq \cdots \leq \frac{\lambda_{k+1}}{C_k} \leq \cdots \leq \frac{\lambda_{n-1}}{C_{n-2}} \leq \frac{\lambda_n}{C_{n-1}} = \frac{1}{n}.$$

The total work of Algorithm 2 depends on k:

 $Work(\mathcal{D}) :=$ cost of 1 iteration number of iterations till ϵ -solution

| k | $P(\mathcal{D})$ | $C(\mathcal{D})$ | $I(\mathcal{D})$ |
|---------------|--|------------------|--|
| 0 | O(n) | O(n) | $\frac{\mathrm{Tr}(\mathbf{A})}{\lambda_1}\ln(1/\epsilon)$ |
| 0 < k < n - 1 | computation of λ_i for $i = 1, 2, \dots, k+1$ computation of u_i for $i = 1, 2, \dots, k$ | O(n) | $\frac{C_k}{\lambda_{k+1}}\ln(1/\epsilon)$ |
| n - 1 | computation of λ_i for $i = 1, 2,, n$ computation of u_i for $i = 1, 2,, n - 1$ | O(n) | $n \ln(1/\epsilon)$ |

6. Bibliography

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