## Stochastic Reformulations

of Linear Systems:

# Algorithms and Convergence Theory 

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Modern Convex Optimization and Applications
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## A System of Linear Equations

$m$ equations with $n$ unknowns


Assumption: The system is consistent (i.e., a solution exists)

## Part I Stochastic Reformulations

P.R. and Martin Takáč

Stochastic Reformulations of Linear Systems: Algorithms and Convergence Theory arXiv:1706.01108, 2017

## Stochastic Reformulations of Linear Systems

## $n \times n$ pos def <br> $B, \mathcal{D}$ <br> $A x=b$

distribution over $m \times q$ matrices

1. Stochastic Optimization
2. Stochastic Linear System
3. Stochastic Fixed Point
4. Probabilistic Intersection

Example: $B=$ identity
$\mathcal{D}=$ uniform over $e_{1}, \ldots, e_{m}\left(\right.$ unit basis vectors in $\left.\mathbb{R}^{m}\right)$

## Theorem

a) These 4 problems have the same solution sets
b) Necessary \& sufficient conditions for the solution set to be equal to $\{x: A x=b\}$

## Reformulation 1: <br> Stochastic Optimization

Minimize $f(x) \stackrel{\text { def }}{=} \mathbf{E}_{S \sim \mathcal{D}}\left[f_{S}(x)\right]$

$$
f_{S}(x)=\frac{1}{2}\left\|x-\Pi_{\mathcal{L}_{S}}^{B}(x)\right\|_{B}^{2}=\frac{1}{2}(A x-b)^{\top} H(A x-b)
$$

$$
\mathcal{L}_{S}=\left\{x: S^{\top} A x=S^{\top} b\right\}
$$

$$
H=S\left(S^{\top} A B^{-1} A^{\top} S\right)^{\dagger} S^{\top}
$$

## Reformulation 2: <br> Stochastic Linear System

Instead of $A x=b$ we solve the preconditioned system:

$$
H=S\left(S^{\top} A B^{-1} A^{\top} S\right)^{\dagger} S^{\top}
$$

$$
\text { Solve } B^{-1} A^{\top} \mathbf{E}_{S \sim \mathcal{D}}[H] A x=B^{-1} A^{\top} \mathbf{E}_{S \sim \mathcal{D}}[H] b
$$

preconditioner
Instead of $B^{-1} A^{\top} \mathbf{E}[H] A$ we have access to $B^{-1} A^{\top} H A$

Unbiased estimate of the preconditioner

## Reformulation 3:

## Stochastic Fixed Point Problem

$$
\text { Solve } x=\mathbf{E}_{S \sim \mathcal{D}}\left[\Pi_{\mathcal{L}_{S}}^{B}(x)\right]
$$

Projection in $B$-norm onto $\mathcal{L}_{S}=\left\{x: S^{\top} A x=S^{\top} b\right\}$

## Reformulation 4: Probabilistic Intersection Problem

Find $x \in \mathbb{R}^{n}$ such that $\mathbf{P}\left(x \in \mathcal{L}_{S}\right)=1$

$$
\mathcal{L}_{S}=\left\{x: S^{\top} A x=S^{\top} b\right\}
$$

Sketched system
$S$ discrete

$$
\left\{x: \mathbf{P}\left(x \in \mathcal{L}_{S}\right)=1\right\}=\bigcap_{S} \mathcal{L}_{S}
$$

## Part II <br> Randomized Algorithms

## Viewpoint 1: Stochastic Optimization

## Stochastic Gradient Descent



A key method in machine learning

## Stochastic "Newton" Descent



## Stochastic Proximal Point Method



# Viewpoint 3: Stochastic Fixed Point Method 

## Stochastic Fixed Point Method

> Stochastic fixed point
> mapping

$$
x^{t+1}=\omega \prod_{\mathcal{L}_{S}}^{B}\left(x^{t}\right)+(1-\omega) x^{t}
$$

Relaxation parameter

$$
S \sim \mathcal{D}
$$

# Part III Complexity 

Basic Method

## Basic Method: Complexity

$$
\mathbf{E}\left[U^{\top} B^{1 / 2}\left(x^{t}-x^{*}\right)\right]=(I-\omega \Lambda)^{t} U^{\top} B^{1 / 2}\left(x^{0}-x^{*}\right)
$$

```
stepsize / relaxation parameter
```

$$
\begin{gathered}
W=B^{-1 / 2} A^{\top} \mathbf{E}_{S \sim \mathcal{D}}[H] A B^{-1 / 2}=U \Lambda U^{\top} \\
H=S\left(S^{\top} A B^{-1} A^{\top} S\right)^{\dagger} S^{\top}
\end{gathered}
$$

## Basic Method: Complexity

Convergence of Expected Iterates
$t \geq \frac{1}{\lambda_{\text {min }}^{+}} \log \left(\frac{1}{\epsilon}\right) \quad \stackrel{\omega=1}{\square}\left\|\mathbf{E}\left[x^{t}-x^{*}\right]\right\|_{B}^{2} \leq \epsilon$
$t \geq \frac{\lambda_{\max }}{\lambda^{+}} \log \left(\frac{1}{\epsilon}\right) \stackrel{\omega=1 / \lambda_{\text {max }}}{\square}\left\|\mathbf{E}\left[x^{t}-x^{*}\right]\right\|_{B}^{2} \leq \epsilon$

L2 Convergence
$t \geq \frac{1}{\lambda_{\text {min }}^{+}} \log \left(\frac{1}{\epsilon}\right) \stackrel{\omega=1}{\longmapsto} \mathbf{E}\left[\left\|x^{t}-x^{*}\right\|_{B}^{2}\right] \leq \epsilon$

Parallel Method

## Parallel Method

"Run 1 step of the basic method from $x^{t}$ several times independently, and average the results."

> i.i.d.

$$
x^{t+1}=\frac{1}{\tau} \sum_{i=1}^{\tau} \phi_{\omega}\left(x^{t}, S_{i}^{t}\right)
$$

One step of the basic method from $x^{t}$

## Parallel Method: Complexity

## L2 Convergence

$$
\begin{array}{cc}
\tau=1 & \tau=+\infty \\
t \geq \frac{1}{\lambda_{\min }^{+}} \log \left(\frac{1}{\epsilon}\right) \quad \text { or } \quad t \geq \frac{\lambda_{\max }}{\lambda_{\min }^{+}} \log \left(\frac{1}{\epsilon}\right)
\end{array}
$$



$$
\mathbf{E}\left[\left\|x^{t}-x^{*}\right\|_{B}^{2}\right] \leq \epsilon
$$

## Accelerated Method

## Accelerated Method

$$
S^{t}, S^{t-1} \sim \mathcal{D} \text { (independent) }
$$

$$
x^{t+1}=\gamma \phi_{\omega}\left(x^{t}, S^{t}\right)+(1-\gamma) \phi_{\omega}\left(x^{t-1}, S^{t-1}\right)
$$

One step of the basic method from $x^{t}$
One step of the basic method from $x^{t-1}$

## Accelerated Method: Complexity

## Convergence of Iterates

$$
t \geq \sqrt{\frac{\lambda_{\max }}{\lambda_{\min }^{+}}} \log \left(\frac{1}{\epsilon}\right) \quad\left\|\mathbf{E}\left[x^{t}-x^{*}\right]\right\|_{B}^{2} \leq \epsilon
$$

$$
\text { Basic Method depends on } \frac{\lambda_{\max }}{\lambda_{\min }^{+}} \text {! }
$$

## Detailed Complexity Results

| Alg. | $\omega$ | $\tau$ | $\gamma$ | Quantity | Rate | Complexity | Theorem |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | - | - | $\left\\|\mathrm{E}\left[x_{k}-x_{*}\right]\right\\|_{\mathbf{B}}^{2}$ | $\left(1-\lambda_{\min }^{+}\right)^{2 k}$ | $1 / \lambda_{\min }^{+}$ | $4.3,4.4,4.6$ |
| 1 | $1 / \lambda_{\max }$ | - | - | $\left\\|\mathrm{E}\left[x_{k}-x_{*}\right]\right\\|_{\mathbf{B}}^{2}$ | $(1-1 / \zeta)^{2 k}$ | $\zeta^{2}$ | $4.3,4.4,4.6$ |
| 1 | $\frac{2}{\lambda_{\min }^{+}+\lambda_{\max }}$ | - | - | $\left\\|\mathrm{E}\left[x_{k}-x_{*}\right]\right\\|_{\mathbf{B}}^{2}$ | $(1-2 /(\zeta+1))^{2 k}$ | $\zeta$ | $4.3,4.4,4.6$ |
| 1 | 1 | - | - | $\mathrm{E}\left[\left\\|x_{k}-x_{*}\right\\|_{\mathbf{B}}^{2}\right]$ | $\left(1-\lambda_{\min }^{+}\right)^{k}$ | $1 / \lambda_{\min }^{+}$ | 4.8 |
| 1 | 1 | - | - | $\mathrm{E}\left[f\left(x_{k}\right)\right]$ | $\left(1-\lambda_{\min }^{+}\right)^{k}$ | $1 / \lambda_{\min }^{+}$ | 4.10 |
| 2 | 1 | $\tau$ | - | $\mathrm{E}\left[\left\\|x_{k}-x_{*}\right\\|_{\mathbf{B}}^{2}\right]$ | $\left(1-\lambda_{\min }^{+}(2-\xi(\tau))\right)^{k}$ |  | 5.1 |
| 2 | $1 / \xi(\tau)$ | $\tau$ | - | $\mathrm{E}\left[\left\\|x_{k}-x_{*}\right\\|_{\mathbf{B}}^{2}\right]$ | $\left(1-\frac{\left.\lambda_{\min }^{+}\right)^{k}}{\xi(\tau)}\right)$ | $\xi(\tau) / \lambda_{\min }^{+}$ | 5.1 |
| 2 | $1 / \lambda_{\max }$ | $\infty$ | - | $\mathrm{E}\left[\left\\|x_{k}-x_{*}\right\\|_{\mathbf{B}}^{2}\right]$ | $(1-1 / \zeta)^{k}$ | $\zeta$ | 5.1 |
| 3 | 1 | - | $\frac{2}{1+\sqrt{0.99 \lambda_{\min }^{+}}}$ | $\left\\|\mathrm{E}\left[x_{k}-x_{*}\right]\right\\|_{\mathbf{B}}^{2}$ | $\left(1-\sqrt{0.99 \lambda_{\min }^{+}}\right)^{2 k}$ | $\sqrt{1 / \lambda_{\min }^{+}}$ | 5.3 |
| 3 | $1 / \lambda_{\max }$ | - | $\frac{2}{1+\sqrt{0.99 / \zeta}}$ | $\left\\|\mathrm{E}\left[x_{k}-x_{*}\right]\right\\|_{\mathbf{B}}^{2}$ | $(1-\sqrt{0.99 / \zeta})^{2 k}$ | $\sqrt{\zeta}$ | 5.3 |

Table 1: Summary of the main complexity results. In all cases, $x_{*}=\Pi_{\mathcal{L}}^{\mathbf{B}}\left(x_{0}\right)$ (the projection of the starting point onto the solution space of the linear system). "Complexity" refers to the number of iterations needed to drive "Quantity" below some error tolerance $\epsilon>0$ (we suppress a $\log (1 / \epsilon)$ factor in all expressions in the "Complexity" column). In the table we use the following expressions: $\xi(\tau)=\frac{1}{\tau}+\left(1-\frac{1}{\tau}\right) \lambda_{\text {max }}$ and $\zeta=\lambda_{\text {max }} / \lambda_{\text {min }}^{+}$.

## Part IV Conclusion

## Contributions

- 4 Equivalent stochastic reformulations of a linear system
- Stochastic optimization
- Stochastic fixed point problem
- Stochastic linear system
- Probabilistic intersection
- 3 Algorithms
- Basic (SGD, stochastic Newton method, stochastic fixed point method, stochastic proximal point method, stochastic projection method, ...)
- Parallel
- Accelerated
- Iteration complexity guarantees for various measures of success
- Expected iterates (closed form)
- L1 / L2 convergence
- Convergence of $f$; ergodic ...


## Related Work

## Basic method with unit stepsize and full rank A:



Robert Mansel Gower and P.R.
Randomized Iterative Methods for Linear Systems
SIAM J. Matrix Analysis \& Applications 36(4):1660-1690, 2015

- 2017 IMA Fox Prize ( $2^{\text {nd }}$ Prize) in Numerical Analysis
- Most downloaded SIMAX paper

Removal of full rank assumption + duality:


Robert Mansel Gower and P.R.
Stochastic Dual Ascent for Solving Linear Systems
arXiv:1512.06890, 2015

Inverting matrices \& connection to Quasi-Newton updates:


Robert Mansel Gower and P.R.
Randomized Quasi-Newton Methods are Linearly Convergent Matrix Inversion Algorithms arXiv:1602.01768, 2016

## Computing the pseudoinverse:



Robert Mansel Gower and P.R.
Linearly Convergent Randomized Iterative Methods for Computing the Pseudoinverse arXiv:1612.06255, 2016

## Application in machine learning:

## THE END

